

CSCI 331:
Introduction to Computer Security

Lecture 7: Password Cracking, part 2

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Topics

Paper discussion (Oechslin)

Precomputed Hash Chains

Rainbow Tables

Your to-dos

1. Project part 1 **due Sunday 9/29.**
2. Reading response (Aleph One), **due Tuesday 10/1.**
3. Lab 2 **due Sunday 10/13.**

Keyed encryption functions

Paper discussion (Oechslin)

Precomputed Hash Chains

Precomputed Hash Chains

Motivation: dictionaries are **too big** to distribute

Recall:



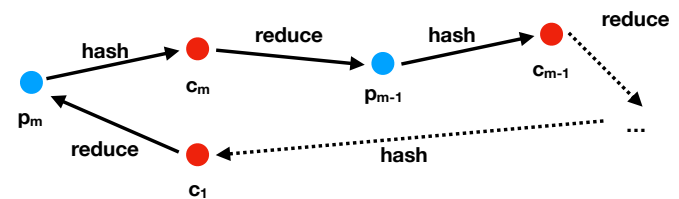
About 26 terabytes!

Thought experiment

- plaintexts
- ciphertexts

Suppose $f(p_i) = c_i$

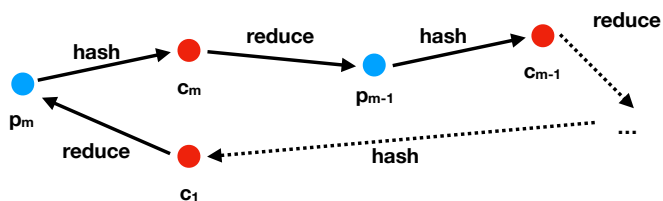
Suppose $r(c_i) = p_{i-1}$ if $i > 1$ otherwise p_m



Thought experiment

● plaintexts
● ciphertexts

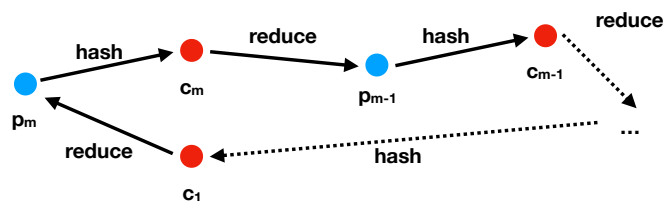
Such a scheme (a *hash chain*) lets us generate all plaintexts (and hashes) from a seed plaintext.



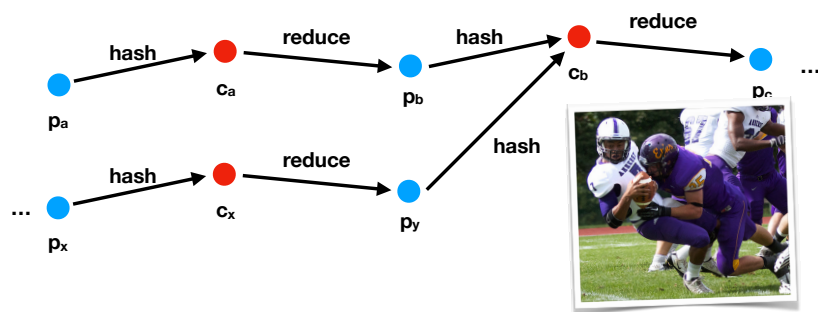
Only need to save the seed. **Drawbacks?**

Thought experiment: **drawbacks**

- Saving just the first password **buys us nothing**. On average, we have to compute $O(m/2)$ hash-reductions to find a password.
- It is **probably not possible** to find a reducer that lets you explore the entire password space.
- **Hash functions collide!**



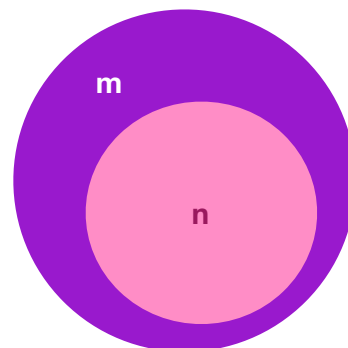
Collisions in a hash chain



After the **collision**, the chain **"loops."**

Collisions prevent us from enumerating the **entire space**!

Hashes are guaranteed to collide

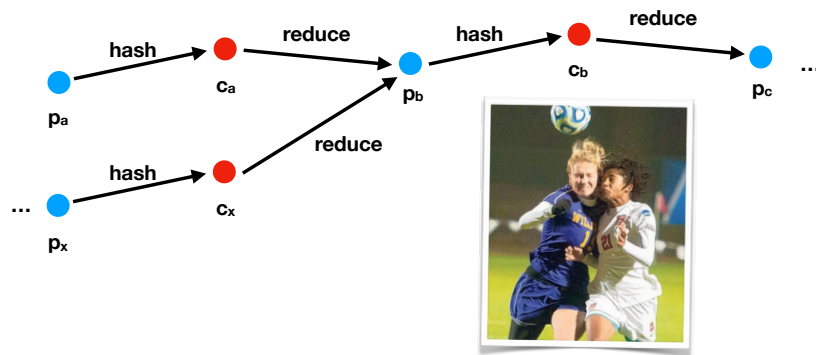


m: # of passwords

n: # of hashes

If $m > n$, we know that **at least** $(m-n)/m$ must collide.

Collisions in a hash chain

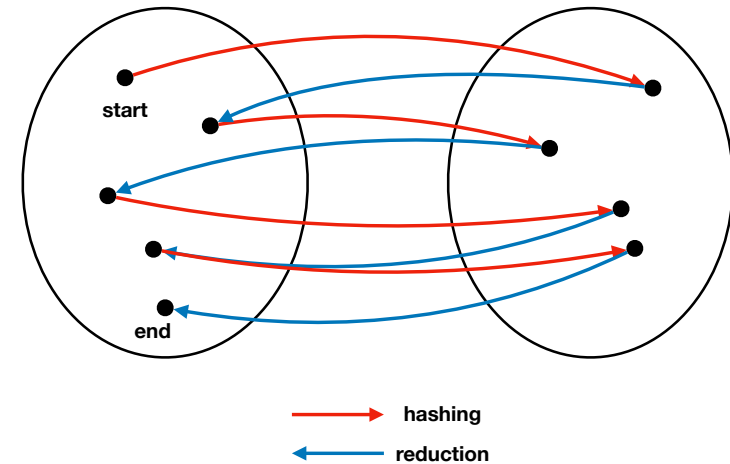


Reducers can produce collisions too!
This is what we mean by an **imperfect reducer**.

Hash chain

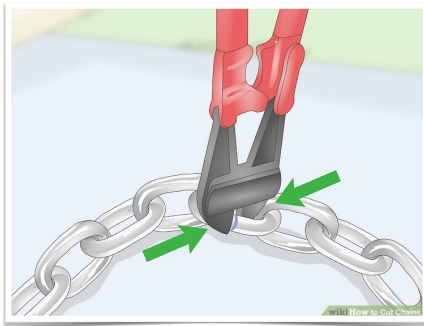
Space of possible passwords

Space of possible hashes



Hash chain of length k

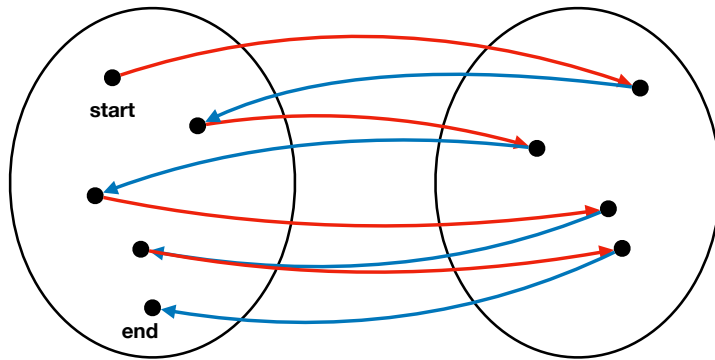
We are going to chop up our long chain into **smaller chains** of length k .



Hash chain

Space of possible passwords

Space of possible hashes

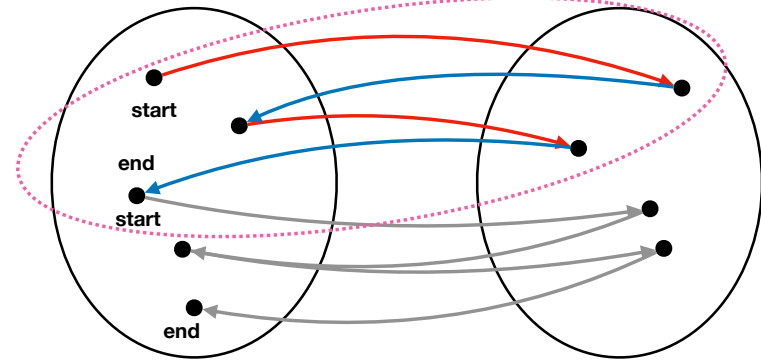


→ hashing
← reduction

Hash chain of length 2

Space of possible passwords

Space of possible hashes

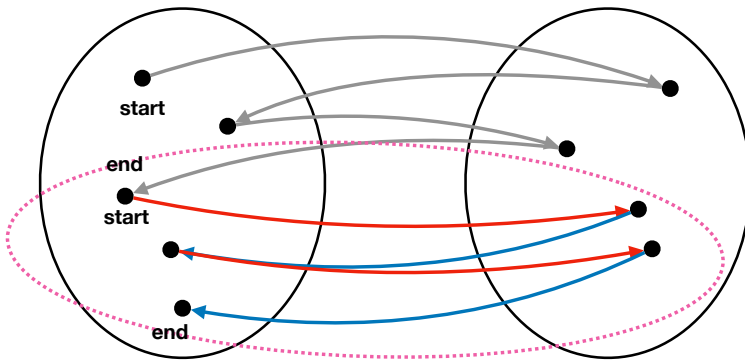


→ hashing
← reduction

Hash chain of length 2

Space of possible passwords

Space of possible hashes



→ hashing
← reduction

Store only start and end

start, end
 p_m , p_{m-3}
 ...
 p_5 , p_3
 p_3 , p_1

Store it backward

end, start
 p_{m-3} , p_m
 ...
 p_3 , p_5
 p_1 , p_3

Suppose you are given c_4

end, start
 p_{m-3} , p_m
 ...
 p_3 , p_5
 p_1 , p_3

reduce
 $c_4 \rightarrow p_3$

Is p_3 an **end point**? **yes**

Hash and **reduce** from **start point**.

password! original ciphertext
 $p_5 \xrightarrow{\text{hash}} c_5 \xrightarrow{\text{reduce}} p_4 \xrightarrow{\text{hash}} c_4$

Hash chain lookup pseudocode

chain table
 ciphertext
 chain length

```
def lookup(c, db, len):
    // look for endpoint
    p = reduce(c)
    i = 0
    while p not end in db && i < len:
        c2 = hash(p)
        p = reduce(c2)
        i++
    if p not end: FAIL
    // we found chain; lookup start pt
    s = db[p]
    c2 = hash(s)
    // decrypt
    i = 0
    while c2 != c && i < len:
        s = reduce(c2)
        c2 = hash(s)
        i++
    if i == len: FAIL
    return s
```

end, start
 ♥♥♥♥♥, ♥♥♥♥♥
 ♥♥♥♥♥, ♥♥♥♥♥
 ♥♥♥♥♥, ♥♥♥♥♥
 ♥♥♥♥♥, ♥♥♥♥♥
 ♥♥♥♥♥, ♥♥♥♥♥

chain table
 Stores **plaintexts**!

Class Activity

Decrypt the hash

7F975A56C761DB6506ECA0B37CE6EC87

Answer:

♥♥♥♥♥

Recap & Next Class

Today we learned:

PCHC algorithm

Next class:

Rainbow algorithm